

Varieties of Cubical Sets

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Abstract. We define a variety of notions of cubical sets, based on sites organized using substructural algebraic theories presenting $\text{PRO}(\mathbf{P})$ s or Lawvere theories. We prove that all our sites are test categories in the sense of Grothendieck, meaning that the corresponding presheaf categories of cubical sets model classical homotopy theory. We delineate exactly which ones are even strict test categories, meaning that products of cubical sets correspond to products of homotopy types.

1 Introduction

There has been substantial interest recently in the use of cubical structure as a basis for higher-dimensional type theories (Angiuli, Harper, and Wilson, 2017; Awodey, 2016; Bezem, Coquand, and Huber, 2014; Cohen, Coquand, Huber, and Mörtberg, 2016). A question that quickly arises is, what sort of cubical structure, because there are several plausible candidates to choose from.

The suitability of a given cubical structure as a basis for a higher-dimensional type theory is dependent on at least two considerations: its proof-theoretic characteristics (e.g. completeness with respect to a given class of models, decidability of equality of terms, existence of canonical forms) and its homotopy-theoretic characteristics, most importantly, that the synthetic homotopy theory to which it gives rise should agree with the standard homotopy theory for topological spaces.

Contributions. In this paper we present a new way to organize the various notions of cubical sets, using substructural algebraic theories as a guiding principle (Mauri, 2005). We define a range of *cube categories* (or *cubical sites*), categories \mathbb{C} for which the corresponding presheaf category $\hat{\mathbb{C}}$ can be thought of as a category of cubical sets. We then consider each of these from the perspective of *test categories* (Grothendieck, 1983), which relates presheaf categories and homotopy theory. We give a full analysis of our cubical sites as test categories.

1.1 Cube Categories

Our cube categories are presented as monoidal categories with a single generating object, X , representing an abstract dimension. To give such a presentation, then, is to give a collection of morphism generators, $f : X^{\otimes n} \rightarrow X^{\otimes m}$, and a collection of equations between parallel morphisms. Such a presentation is known in the literature as a “PRO”. Two closely-related notions are those of “PROP” and of “Lawvere theory”, the difference being that a PROP assumes that the monoidal category is *symmetric*, and a Lawvere theory further assumes that it is *cartesian*. These additional assumptions manifest themselves in the proof theory as the admissibility of certain structural rules: in a PROP, the structural rule of *exchange* is admissible, and in a Lawvere theory *weakening* and *contraction* are admissible as well.

The distinguishing property of a theory making it “cubical”, rather than of some other “shape”, is the presence of two generating morphisms, $d^0, d^1 : 1 \rightarrow X$ representing the face maps (where “1” is the monoidal unit). This is because combinatorially, an n -dimensional cube has $2n$ many $(n - 1)$ -cube faces, namely, two in each dimension.

1.2 Test Categories

In his epistolary research diary, *Pursuing Stacks*, Grothendieck set out a program for the study of abstract homotopy. The homotopy theory of topological spaces can be described as that of (weak) higher-dimensional groupoids. Higher-dimensional groupoids also arise from combinatorially-presented structures, such as simplicial or, indeed, cubical sets, which can thus be seen as alternative presentations of the classical homotopy category, **Hot**. That is, **Hot** arises either from the category of topological spaces, or from the category of simplicial sets, by inverting a class of morphisms (the weak equivalences).

Recall that for any small category A and any cocomplete category \mathcal{E} , there is an equivalence between functors $i : A \rightarrow \mathcal{E}$ and adjunctions

$$\widehat{A} \begin{array}{c} \xleftarrow{i^*} \\ \xrightarrow[i_!]{\tau} \end{array} \mathcal{E} \quad (1)$$

where $i_!$ is the left Kan extension of i , and $i^*(X)(a) = \text{Hom}_{\mathcal{E}}(i(a), X)$.³ Here \widehat{A} is the category of functors $A^{\text{op}} \rightarrow \mathbf{Set}$, also known as the presheaf topos of A . The category of simplicial sets is defined as the presheaf category $\mathbf{sSet} = \widehat{\Delta}$, where Δ is the category of non-empty finite totally ordered sets and order-preserving maps.

³ For readers familiar with the new approach to categories in homotopy type theory (Univalent Foundations), as in Ahrens, Kapulkin, and Shulman (2015), we emphasize that A is a *strict* category, that is, a precategory based on a set of objects. In this paper, all small categories are considered as strict categories, and **Cat** denotes the 1-category of these.

The category of small categories, \mathbf{Cat} , likewise presents the homotopy category, by inverting the functors that become weak equivalences of simplicial sets after applying the nerve functor, which is the right adjoint functor $N = i^*$ coming from the inclusion $\Delta \hookrightarrow \mathbf{Cat}$. Grothendieck realized that this allows us to compare the presheaf category \widehat{A} , for any small category A , with the homotopy category in a canonical way, via the functor $i_A : A \rightarrow \mathbf{Cat}$ given by $i_A(a) = A/a$. The left Kan extension is in this case also denoted i_A and we have $i_A(X) = A/X$ for X in \widehat{A} . He identified the notion of a *test category*, which is one for which the homotopy category of its category of presheaves is equivalent to the homotopy category via the right adjoint i_A^* , after localization at the weak equivalences. In other words, presheaves over a test category are models for homotopy types of spaces/ ∞ -groupoids. In this way, the simplex category Δ is a test category, and the classical cube category, $\mathbb{C}_{(w, \cdot)}$ in our notation, is also a test category.

One perceived benefit of the simplex category Δ is that the induced functor $\mathbf{sSet} \rightarrow \mathbf{Hot}$ preserves products, whereas the functor $\widehat{\mathbb{C}_{(w, \cdot)}} \rightarrow \mathbf{Hot}$ does not. A test category A for which the functor $\widehat{A} \rightarrow \mathbf{Hot}$ preserves products is called a *strict test category*. We shall show that most natural cube categories are in fact strict test categories.

2 Cube categories

In this paper, we consider as base categories the syntactic categories for a range of monoidal theories, all capturing some aspect of the notion of an interval. We call the resulting categories “cube categories”, because the monoidal powers of the interval then correspond to cubes⁴.

The generating morphisms for these cube categories can be classified either as *structural*, natural families corresponding to structural rules of a proof theory, or as *algebraic*, distinguished by the property of having coarity one. Our cube categories vary along three dimensions: the structural rules present, the signature of algebraic function symbols, and the equational theory.

This section is organized as follows: in subsection 2.1 we recall the basics of algebraic theories in monoidal categories, and in subsection 2.2 we discuss how monoidal languages are interpreted in monoidal categories. Then we introduce in subsection 2.3 the monoidal languages underlying our cubical theories, so that in subsection 2.4 we can introduce the standard interpretations of these and our cubical theories. In subsection 2.5 we give a tour of the resulting cube categories.

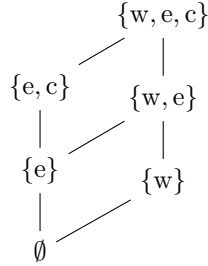
2.1 Monoidal algebraic theories

We assume the reader is familiar with ordinary algebraic theories and their categorical incarnations as Lawvere theories. The idea of monoidal algebraic theories as described by Mauri (2005) is to generalize this to cases where only a

⁴ Here we say “cube” for short instead of “hypercube” for an arbitrary dimensional power of an interval. A prefix will indicate the dimension, as in 0-cube, 1-cube, etc.

subset of the structural rules (weakening, exchange and contraction) are needed to describe the axioms. Think for example of the theory of monoids over a signature of a neutral element 1 and a binary operation. The axioms state that $1x = x = x1$ (in the context of one variable x) and $x(yz) = (xy)z$ (in the context of the variables x, y and z). Note that the terms in these equations each contain all the variables of the context and in the same order. Thus, no structural rules are needed to form the equations, and hence the notion of a monoid makes sense in any monoidal category.

Let us now make this more precise. For the structural rules, we follow Mauri (*ibid.*) and consider any subset of $\{w, e, c\}$ (w for weakening, e for exchange, and c for contraction), except that whenever c is present, so is e . Thus we consider the following lattice of subsets of structural rules:



It would be possible to consider monoidal theories over operations that have any number of incoming and outgoing edges, representing morphisms $f : A_1 \otimes \cdots \otimes A_n \rightarrow B_1 \otimes \cdots \otimes B_m$ in a monoidal category. However, it simplifies matters considerably, and suffices for our purposes, to consider only operations of the usual kind in algebra, with any arity of incoming edges and exactly one outgoing edge, as in $f : A_1 \otimes \cdots \otimes A_n \rightarrow B$.

Furthermore, we shall consider only single-sorted theories, defined over a signature Σ given by a set of function symbols with arities. For a single-sorted signature, the types can of course be identified with the natural numbers.

A (single-sorted, algebraic) *monoidal language* consists of a pair of a set of structural rules together with an algebraic signature. A term t in a context of n free variables x_1, \dots, x_n is built up from the function symbols in such a way that when we list the free variables in t from left to right:

- every variables occurs unless w (weakening) is a structural rule;
- the variables occur in order unless e (exchange) is a structural rule; and
- there are no duplicated variables unless c (contraction) is a structural rule.

For the precise rules regarding term formation and the proof theory of equalities of terms, we refer to Mauri (2005).

2.2 Interpretations

An interpretation of a monoidal language in a monoidal category $(\mathcal{E}, 1, \otimes)$ will consist of an object X representing the single sort, together with morphisms

representing the structural rules and morphisms representing the function symbols (including constants).

The structural rules w , e , c are interpreted respectively by morphisms

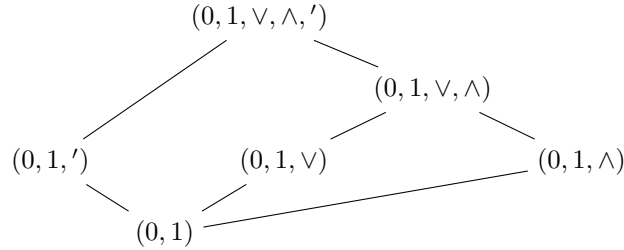
$$\begin{aligned}\varepsilon &: X \rightarrow 1 \\ \tau &: X \otimes X \rightarrow X \otimes X \\ \delta &: X \rightarrow X \otimes X\end{aligned}$$

satisfying certain laws (Mauri, 2005, (27–36)). These laws specify also the interaction between these morphisms and the morphisms corresponding to the function symbols. When \mathcal{E} is symmetric monoidal, we interpret τ by the braiding of \mathcal{E} , and when \mathcal{E} is cartesian monoidal, we interpret everything using the cartesian structure. A function symbol f of arity n is interpreted by a morphism $|f| : X^{\otimes n} \rightarrow X$. This morphism and the structural morphisms are required to interact nicely, e.g., $\varepsilon \circ |f| = \varepsilon^{\otimes n}$ (cf. *loc. cit.*).

In a syntactic category for the empty theory (of which the syntactic category for a non-empty theory is a quotient), every morphism factors uniquely as a structural morphism followed by a functional one (*op. cit.*, Prop. 5.1). When we impose a theory, we may of course lose uniqueness, but we still have existence.

2.3 Cubical monoidal languages

All of our cubical signatures will include the two endpoints of the interval as nullary function symbols, 0 and 1 . For the rest, we consider the two “connections” \vee and \wedge , as well as the reversal, indicated by a prime, $'$. This gives us the following lattice of 6 signatures:



Combined with the 6 possible combinations of structural rules, we are thus dealing with 36 distinct languages.

Definition 1. Let $L_{(a,b)}$, where a is a substring of “wec” and b a substring of “ $\vee \wedge '$ ”, denote the language with structure rules from a and signature obtained from $(0, 1)$ by expansion with the elements of b .

We shall, however, mostly consider the 18 languages with weakening present (otherwise we are dealing with *semi-cubical sets*).

The third dimension of variation for our cube categories is the algebraic theory for a given language. Here we shall mostly pick the theory of a particular standard structure, so let us pause our discussion of cube categories to introduce our standard structures.

2.4 The canonical cube categories

We consider the interval structures of the following objects in the cartesian monoidal categories **Top** and **Set** (so all structural rules are supported, interpreted by the global structure):

- the standard topological interval, $I = [0, 1]$ in **Top**;
- the standard 2-element set, $I = 2 = \{0, 1\}$ in **Set**.

For the topological interval, we take $x \vee y = \max\{x, y\}$, $x \wedge y = \min\{x, y\}$ and $x' = 1 - x$. These formulas also apply to the 2-element set 2.

We can also consider the 3-element Kleene algebra (cf. subsection 2.5) $\mathbb{3} = \{0, u, 1\}$ with $u = u'$ and the 4-element de Morgan algebra $\mathbb{D} = \{0, u, v, 1\}$ with $u' = u$ and $v' = v$ in **Set** (called the *diamond*), as further structures. It is interesting to note here that $\mathbb{3}$ has the same theory as $[0, 1]$ in the full language, namely the theory of Kleene algebras (Gehrke, Walker, and Walker, 2003). Of course, 2 gives us the theory of Boolean algebras, while \mathbb{D} gives the theory of de Morgan algebras.

Definition 2. With (a, b) as in Def. 1, and T a theory in the language $L_{(a,b)}$, let $\mathbb{C}_{(a,b)}(T)$ denote the syntactic category of the theory T . We write $\mathbb{C}_{(a,b)}$ for short for $\mathbb{C}_{(a,b)}(\text{Th}([0, 1]))$ for the syntactic category of the topological interval with respect to the monoidal language $L_{(a,b)}$.

We say that $\mathbb{C}_{(a,b)}$ is the *canonical* cube category for the language $L_{(a,b)}$.

Proposition 1. The canonical cube category $\mathbb{C}_{(a,b)}$ is isomorphic to the monoidal subcategory of **Top** generated by $[0, 1]$ with respect to the language $L_{(a,b)}$.

In fact, since the forgetful functor **Top** \rightarrow **Set** is faithful, we get the same theory in every language whether we consider the interval $[0, 1]$ in **Top** or in **Set**, and it generates the same monoidal subcategory in either case.

Proposition 2. If L is a cubical language strictly smaller than the maximal language, $L_{(\text{wec}, \vee \wedge')}$, then the theory of the standard structure $[0, 1]$ is the same as the theory of the standard structure $2 = \{0, 1\}$ in **Set**. For $L_{(\text{wec}, \vee \wedge')}$ the theory of $[0, 1]$ equals the theory of the three-element Kleene algebra, $\mathbb{3}$, and is in fact the theory of Kleene algebras.

Proof. If L does not have reversal, then we can find for any points $a < b$ in $[0, 1]$ a homomorphism $f : [0, 1] \rightarrow 2$ in **Set** with $f(a) < f(b)$, so $[0, 1]$ and 2 have the same theory.

If L does not have contraction, each term is monotone in each variable. By pushing reversals towards the variables and applying the absorption laws, we can find for each resulting term $s[x, y_1, \dots, y_n]$ instantiations $b_1, \dots, b_n \in 2$ with $s[0, b_1, \dots, b_n] \neq s[1, b_1, \dots, b_n]$. It follows that if two terms agree on instantiations in 2, then they use each variable the same way: either with or without a reversal. Hence we can ignore the reversals and use the previous case.

The last assertion is proved by Gehrke et al. (2003). □

Axiom	Lang. req.	Name
$x \vee (y \vee z) = (x \vee y) \vee z$	(\cdot, \vee)	\vee -associativity
$0 \vee x = x = x \vee 0$	(\cdot, \vee)	\vee -unit
$1 \vee x = 1 = x \vee 1$	(w, \vee)	\vee -absorption
$x \vee y = y \vee x$	(e, \vee)	\vee -symmetry
$x \vee x = x$	(ec, \vee)	\vee -idempotence
$x \wedge (y \wedge z) = (x \wedge y) \wedge z$	(\cdot, \wedge)	\wedge -associativity
$1 \wedge x = x = x \wedge 1$	(\cdot, \wedge)	\wedge -unit
$0 \wedge x = 0 = x \wedge 0$	(w, \wedge)	\wedge -absorption
$x \wedge y = y \wedge x$	(e, \wedge)	\wedge -symmetry
$x \wedge x = x$	(ec, \wedge)	\wedge -idempotence
$x'' = x$	$(\cdot, ')$	$'$ -involution
$0' = 1$	$(\cdot, ')$	$'$ -computation
$x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$	$(ec, \vee \wedge)$	distributive law 1
$x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$	$(ec, \vee \wedge)$	distributive law 2
$x = x \vee (x \wedge y) = x \wedge (x \vee y)$	$(wec, \vee \wedge)$	lattice-absorption
$(x \vee y)' = x' \wedge y'$	$(\cdot, \vee \wedge')$	de Morgan's law
$x \wedge x' \leq y \vee y'$	$(wec, \vee \wedge')$	Kleene's law

Table 1. Cubical axioms

Corollary 1. *For each canonical cube category $\mathbb{C}_{(a,b)}$ we have decidable equality of terms.*

Proof. This follows because equality is decidable in the theory of Kleene algebras, as this theory is characterized by the object $\mathfrak{3}$. Hence we can decide equality by the method of “truth-tables”, however, relative to the elements of $\mathfrak{3}$. \square

Of course, any algorithm for equality in the theory of Kleene algebras can be used, including the one based on disjunctive normal forms (*op. cit.*).

We can in the presence of weakening quite simply give explicit axiomatizations for each of the canonical cube categories $\mathbb{C}_{(a,b)}$: take the subset of axioms in Tab. 1 that make sense in the language $L_{(a,b)}$. Without weakening, this will not suffice, as we should need infinitely many axioms to ensure for instance that $s \wedge 0 = t \wedge 0$ where s, t are terms in the same variable context.

When we compare our canonical cube categories to the categories of Grandis and Mauri (2003), we find $\mathbb{I} = \mathbb{C}_{(w, \cdot)}$, $\mathbb{J} = \mathbb{C}_{(w, \vee \wedge)}$, $\mathbb{K} = \mathbb{C}_{(we, \vee \wedge)}$, and $!\mathbb{K} = \mathbb{C}_{(we, \vee \wedge')}$.

For the full language, $L_{(wec, \vee \wedge')}$, there are two additional interesting cube categories: \mathbb{C}_{dM} (de Morgan algebras) and \mathbb{C}_{BA} (boolean algebras), where de Morgan algebras satisfy all the laws of Tab. 1 except Kleene's law, and boolean algebras of course satisfy additionally the law $x \vee x' = 1$. The case of de Morgan algebras is noteworthy as the basis of the cubical model of type theory of Cohen et al. (2016).

Definition 3. For each of our cube categories \mathbb{C} , we write $[n]$ for the object representing a context of n variables, and we write \square^n for the image of $[n]$ under the Yoneda embedding $y : \mathbb{C} \hookrightarrow \widehat{\mathbb{C}}$.

2.5 A tour of the menagerie

Plain cubes In the canonical cube category for the language $L_{(w, \cdot)}$ the monoidal unit $[0]$ is terminal and the interpretation of weakening is generated by the unique degeneracy map $\varepsilon : [1] \rightarrow [0]$. The points 0 and 1 are interpreted by face maps $\eta_0, \eta_1 : [0] \rightarrow [1]$. In the following we let $\{i, j\}$ range over $\{0, 1\}$ with the assumption that $i \neq j$.

Because $[0]$ is terminal we have the *face-degeneracy laws*, $\varepsilon \circ \eta_i = \text{Id}_{[0]}$, which we represent string-diagrammatically as:

$$\boxed{\begin{array}{c} \eta_i \\ \varepsilon \end{array}} = \boxed{}$$

This law is invisible in the algebraic notation.

In a presheaf T , the face maps give rise to reindexing functions between the fibers, $\partial_k^i : T[n+1] \rightarrow T[n]$, where $1 \leq k \leq n+1$. In particular, when n is 0 these pick out the respective boundary points of an interval. Similarly, the degeneracy map gives rise to reindexing functions between the fibers, $*_k : T[n] \rightarrow T[n+1]$, where $1 \leq k \leq n+1$. In particular, when n is 0 this determines a degenerate interval $*(a) \in T[1]$ on a point $a \in T[0]$, and the face-degeneracy laws tell us its boundary points: $\partial^0(*(a)) = \partial^1(*(a)) = a$.

Adding the structural law of exchange to the language adds a natural isomorphism $\tau : [2] \rightarrow [2]$ to the syntactic category. In a presheaf, this lets us permute any adjacent pair of a cube's dimensions by reflecting it across the corresponding diagonal (hyper)plane. In particular, $T(\tau)$ reflects a square across its main diagonal.

Cubes with diagonals In the canonical cube category for the language $L_{(wec, \cdot)}$ the monoidal product is cartesian and the interpretation of contraction is generated by the diagonal map $\delta : [1] \rightarrow [2]$. In this case, the pair (δ, ε) forms a comonoid.

Because δ is natural we have the *face-diagonal laws*, $\delta \circ \eta_i = \eta_i \otimes \eta_i$, which we represent string-diagrammatically as:

$$\boxed{\begin{array}{c} \eta_i \\ \delta \end{array}} = \boxed{\begin{array}{cc} \eta_i & \eta_i \end{array}}$$

In a presheaf T , the diagonal map gives rise to reindexing functions between the fibers, $d_k : T[n+2] \rightarrow T[n+1]$, where $1 \leq k \leq n+1$. In particular, when n is 0 this picks out the main diagonal of a square, and the face-diagonal laws tell us its boundary points.

In the cubical semantics, the fact that ε is a comonoid counit for δ tells us that the diagonal of a square formed by degenerating an interval is just that interval. Likewise, the fact that δ is coassociative tells us that the main diagonal interval of a higher-dimensional cube is well-defined.

Cubes with reversals In the canonical cube category for the language $L_{(\cdot, \cdot)}$ we have an involutive reversal map $\rho : [1] \rightarrow [1]$. A reversal acts by swapping the face maps, giving the *face-reversal laws*, $\rho \circ \eta_i = \eta_j$, which we represent string-diagrammatically as:

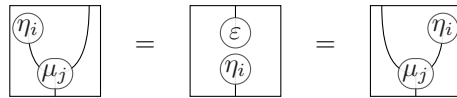


This embodies the equations $0' = 1$ and $1' = 0$.

In a presheaf T , the reversal map gives rise to endomorphisms on the fibers, $!_k : T[n+1] \rightarrow T[n+1]$, where $1 \leq k \leq n+1$. In particular, when n is 0 this reverses an interval, and the face-reversal laws tell us its boundary points.

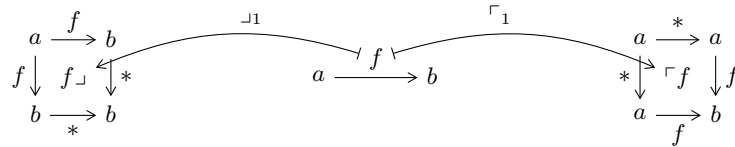
If the signature contains weakening, then we have the *reversal-degeneracy law*, $\varepsilon \circ \rho = \varepsilon$. In the cubical interpretation, this says that a degenerate interval is invariant under reversal. If the signature contains contraction, then we have the *reversal-diagonal law*, $\delta \circ \rho = (\rho \otimes \rho) \circ \delta$. Cubically, this says that the reversal of a square's diagonal is the diagonal of the square resulting from reversing in each dimension.

Cubes with connections In the canonical cube category for the language $L_{(w, \vee \wedge)}$, the connectives are interpreted by connection maps $\mu_0, \mu_1 : [2] \rightarrow [1]$. Each pair (μ_i, η_i) forms a monoid. Furthermore, the unit for each monoid is an absorbing element for the other: $\mu_j \circ (\eta_i \otimes \text{Id}) = \eta_i \circ \varepsilon = \mu_j \circ (\text{Id} \otimes \eta_i)$. Grandis and Mauri (2003) call such structures “dioids”, so we refer to these as the *dioid absorption laws*, and represent them string-diagrammatically as:



Equationally, these amount to the \vee - and \wedge -absorption laws of Tab. 1.

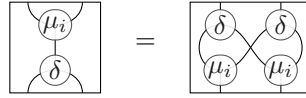
In the cubical semantics, connections can be seen as a variant form of degeneracy, identifying adjacent, rather than opposite, faces of a cube. Of course in order for this to make sense, a cube must have at least two dimensions. In a presheaf T , the connection maps give rise to reindexing functions between the fibers, $\ulcorner_k, \lrcorner_k : T[n+1] \rightarrow T[n+2]$, where $1 \leq k \leq n+1$. In particular, when n is 0 these act as follows:



It may help to think of the interval f as a folded paper fan, with its “hinge” at the domain end in the case of \lrcorner , and at the codomain end in the case of \lrcorner . The respective connected squares are then obtained by “opening the fan”. The monoid unit laws give us the (generally) non-degenerate faces of a connected square, while the dioid absorption laws give us the (necessarily) degenerate ones. Monoid associativity says that multiply-connected higher-dimensional cubes are well-defined.

Because $[0]$ is terminal we also have the *connection-degeneracy laws*, $\varepsilon \circ \mu_i = \varepsilon \otimes \varepsilon$, which say that a connected square arising from a degenerate interval is a doubly-degenerate square. In the presence of exchange, we further assume that the μ_i are commutative. In the cubical semantics, this implies that connected cubes are invariant under reflection across their connected dimensions.

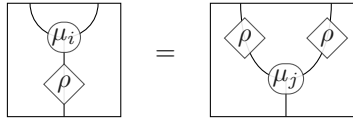
When adding contraction to the signature, we must give a law for rewriting a diagonal following a connection such that the structural maps come first. This is done by the *connection-diagonal laws*, $\delta \circ \mu_i = (\mu_i \otimes \mu_i) \circ (\text{Id} \otimes \tau \otimes \text{Id}) \circ (\delta \otimes \delta)$, which we represent string-diagrammatically as:



Algebraically, this says that two copies of a conjunction consists of a pair of conjunctions, each on copies of the respective terms; cubically, it gives the connected square on the diagonal interval of another square in terms of a product of diagonals in the 4-cube resulting from connecting each dimension separately. The whole structure, then, is a pair of bicommutative bimonoids, $(\delta, \varepsilon, \mu_i, \eta_i)$ related by the dioid absorption laws. We refer to this as a *linked bimonoid*.

Additionally, we may impose the laws of bounded, modular, and distributive lattices, each of which implies its predecessors, and all of which imply the *diagonal-connection laws*, $\mu_i \circ \delta = \text{Id}_{[1]}$, known in the literature (rather generically) as “special” laws. These correspond to the \vee - and \wedge -idempotence laws of Tab. 1.

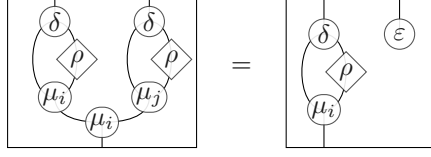
The full signature When both reversals and connections are present, the *de Morgan laws*, $\rho \circ \mu_i = \mu_j \circ (\rho \otimes \rho)$ permute reversals before connections:



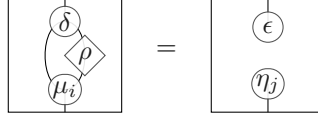
These laws imply each other as well as their nullary versions, the face-reversal laws.

Using the algebraic characterization of order in a lattice, $x \wedge y = x \iff x \leq y \iff x \vee y = y$, we can express the *Kleene law* as $\mu_i \circ (\mu_i \otimes \mu_j) \circ (\text{Id} \otimes \rho \otimes \rho)$

$$\text{Id} \otimes \rho) \circ (\delta \otimes \delta) = \mu_i \circ (\text{Id} \otimes \rho) \circ (\delta \otimes \varepsilon):$$



Finally, we arrive at the structure of a boolean algebra by assuming the *Hopf laws*, $\mu_i \circ (\text{Id} \otimes \rho) \circ \delta = \eta_j \circ \varepsilon$:



Cubically, these say that the anti-diagonal of a connected square is degenerate.

3 Test categories

As mentioned in the introduction, a test category is a small category A such that the presheaf category \hat{A} can model the homotopy category after a canonical localization. Here we recall the precise definitions (cf. Maltsiniotis (2005) for more details).

The inclusion functor $i : \Delta \rightarrow \mathbf{Cat}$ induces via (1) an adjunction $i_! : \hat{\Delta} \rightleftarrows \mathbf{Cat} : N$ where the right adjoint N is the *nerve* functor. This allows us to transfer the homotopy theory of simplicial sets to the setting of (strict) small categories, in that we define a functor $f : A \rightarrow B$ to be a *weak equivalence* if $N(f)$ is a weak equivalence of simplicial sets. A small category A is called *aspheric* (or *weakly contractible*) if the canonical functor $A \rightarrow 1$ is a weak equivalence. Since any natural transformation of functors induces a homotopy, it follows that any category with a natural transformation between the identity functor and a constant endofunctor is contractible, and hence aspheric. In particular, a category with an initial or terminal object is aspheric. It follows from Quillen's Theorem A that a functor $f : A \rightarrow B$ is a weak equivalence, if it is *aspheric*, meaning that all the slice categories $A_{/b}$ are aspheric, for $b \in B$.⁵ By duality, so is any *coaspheric* functor $f : A \rightarrow B$, one for which the coslice categories ${}_b \backslash A$, for $b \in B$, are all aspheric. Finally, we say that a presheaf X in \hat{A} is *aspheric*, if $A_{/X}$ is.

For any small category A , we can use the adjunction induced by the functor $i_A : A \rightarrow \mathbf{Cat}$, $i_A(a) = A_{/a}$ (as mentioned in subsection 1.2) to define the class of weak equivalences in \hat{A} , \mathcal{W}_A : $f : X \rightarrow Y$ in \hat{A} is in \mathcal{W}_A if $i_A(f) : A_{/X} \rightarrow A_{/Y}$ is a weak equivalence of categories. We make the following definitions:

- A is a *weak test category* if the induced functor $\overline{i_A} : (\mathcal{W}_A)^{-1} \hat{A} \rightarrow \mathbf{Hot}$ is an equivalence of categories.

⁵ The slice $A_{/b}$ denotes the comma category $(f \downarrow b)$ corresponding to the functors $f : A \rightarrow B$ and $b : 1 \rightarrow B$. Similarly, the coslice ${}_b \backslash A$ denotes $(b \downarrow f)$.

$a \setminus b$	\cdot	$'$	\vee	\wedge	$\vee \wedge$	$\vee \wedge'$
w	t	t	st	st	st	st
we	t	t	st	st	st	st
wec	st	st	st	st	st	st/st/st

Table 2. Which canonical cube categories $\mathbb{C}_{(a,b)}$ are test (t) or even strict test (st) categories. The bottom-right corner refers to the cube categories corresponding to de Morgan, Kleene and boolean algebras.

- A is a *local test category* if all the slices A/a are weak test categories.
- A is a *test category* if A is both a weak and a local test category.
- A is a *strict test category* if A is a test category and the functor $\hat{A} \rightarrow \mathbf{Hot}$ preserves finite products.

The goal of the rest of this section is to prove that all the cube categories are test categories, and to establish exactly which ones are strict test categories. First we introduce Grothendieck interval objects. These allow us to show that any cartesian cube category is a strict test category (Cor. 2) as well as to show that all the cube categories are test categories (Cor. 3). Then we adapt an argument of Maltsiniotis (2009) to show that any cube category with a connection is a strict test category (Thm. 3). Finally, we adapt another argument of his to show that the four remaining cube categories fail to be strict test categories (Thm. 4). The end result is summarized in Tab. 2.

In order to study test categories, Grothendieck (1983) introduced the notion of an *interval* (*segment* in the terminology of Maltsiniotis (2005)) in a presheaf category \hat{A} . This is an object I equipped with two global elements d^0, d^1 . As such, it is just a structure for the initial cubical language, $L_{(\cdot, \cdot)}$ in the cartesian monoidal category \hat{A} , and \square^1 is thus canonically a Grothendieck interval in all our categories of cubical sets, $\hat{\mathbb{C}}$.

An interval is *separated* if the equalizer of d^0 and d^1 is the initial presheaf, \emptyset . For cubical sets, this is the case if and only if $0 = 1$ is not derivable in the base category, in any variable context.

The following theorem is due to Grothendieck (1983, 44(c)), cf. Thm. 2.6 of Maltsiniotis (2009). We say that A is *totally aspheric* if A is non-empty and all the products $y(a) \times y(b)$ are aspheric, where $y : A \rightarrow \hat{A}$ is the Yoneda embedding.

Theorem 1 (Grothendieck). *If A is a small category that is totally aspheric and has a separated aspheric interval, then A is a strict test category.*

Corollary 2. *If \mathbb{C} is any canonical cube category over the full set of structural rules, $\mathbb{C}_{(\text{wec}, b)}$, or the cartesian cube category of de Morgan algebras or that of boolean algebras, then \mathbb{C} is a strict test category.*

Proof. Since \mathbb{C} has finite products it is totally aspheric, as the Yoneda embedding preserves finite products. The Grothendieck interval corresponding to the 1-cube

is representable and hence aspheric. This interval is separated as $0 = 1$ is not derivable in any context. \square

The following theorem is contained in the Proposition of Grothendieck (1983, 44(d), p. 86):

Theorem 2 (Grothendieck). *If A is a small aspheric category with a separated aspheric interval (I, δ_0, δ_1) in \widehat{A} , and $i : A \rightarrow \mathbf{Cat}$ is a functor such that for any a in A , $i(a)$ has a final object, and there is a map of intervals $i_!(I) \rightarrow 2$ in \mathbf{Cat} , then A is a test category.*

In this case, i is in fact a *weak test functor*, meaning that $i^* : \mathbf{Cat} \rightarrow \widehat{A}$ induces an equivalence $\mathbf{Hot} \rightarrow (\mathcal{W}_A)^{-1}\widehat{A}$.

Corollary 3. *Any canonical cube category \mathbb{C} is a test category.*

Proof. The conditions of the theorem hold trivially for any cube category without reversal, taking i to be the canonical functor sending I to 2. If we have reversals, then we can define i by sending I to the category with three objects, $\{0\}, \{1\}, \{0, 1\}$ and arrows $\{0\}, \{1\} \rightarrow \{0, 1\}$. Since this is a Kleene algebra, this functor is well-defined in all cases. \square

The cube categories \mathbb{C}_{dM} and \mathbb{C}_{BA} of de Morgan and boolean algebras are test categories by Cor. 2.

We now turn to the question of which non-cartesian cube categories are strict test categories. The following theorem was proved by Maltsiniotis (2009, Prop. 3.3) for the case of $\mathbb{C}_{(\text{w}, \vee)}$. The same proof works more generally, so we obtain:

Theorem 3. *Any canonical cube category $\mathbb{C}_{(a, b)}$ where b includes one of the connections \vee, \wedge , is a strict test category.*

That leaves four cases: (w, \cdot) , $(\text{w}, ')$, (we, \cdot) , $(\text{we}, ')$. The first of these, the “classical” cube category, is not a strict test category by the argument of Maltsiniotis (2009, Sec. 5).

Note the unique factorizations we have for these categories: every morphism $f : [m] \rightarrow [n]$ factors as degeneracies, followed by (possibly) symmetries, followed by (possibly) reversals, followed by face maps. The isomorphisms are exactly the compositions of reversals and symmetries (if any).

Next, we use variations of the argument of Maltsiniotis (*loc. cit.*) to show that none of these four sites are strict test categories by analyzing the homotopy type of the slice category $\mathbb{C}_{/\square^1 \times \square^1}$ in each case.

Theorem 4. *The canonical cube categories $\mathbb{C}_{(\text{w}, \cdot)}$, $\mathbb{C}_{(\text{w}, ')}$, $\mathbb{C}_{(\text{we}, \cdot)}$ and $\mathbb{C}_{(\text{we}, ')}$ are not strict test categories.*

Proof. Let \mathbb{C} be one of these categories. We shall find a full subcategory A of $\mathbb{C}_{/\square^1 \times \square^1}$ that is not aspheric, and such that the inclusion is a weak equivalence. Hence $\mathbb{C}_{/\square^1 \times \square^1}$ is not aspheric, and \mathbb{C} cannot be a strict test category.

An object of the slice category $\mathbb{C}/\square^1 \times \square^1$ is given by a dimension and two terms corresponding to that dimension. We can thus represent it by a variable context $x_1 \cdots x_n$ (for an n -cube $[n]$) and two terms s, t in that context.

For the classical cube category case (w, \cdot) , we let A contain the following objects, cf. Maltiniotis (2009, Prop. 5.2):

- 4 zero-dimensional objects $(\cdot, (i, j))$, $i, j \in \{0, 1\}$.
- 5 one-dimensional objects; 4 corresponding to the sides of a square, $(x, (i, x))$ and $(x, (x, i))$ with $i \in \{0, 1\}$, and 1 corresponding to its diagonal, $(x, (x, x))$.
- 2 two-dimensional objects: $(xy, (x, y))$ and $(xy, (y, x))$ (let us call them, respectively, the northern and southern hemispheres).

In the presence of the exchange rule, we do not need both of the two-dimensional objects (so we retain, say, the northern hemisphere), and in the presence of reversal, we need additionally the anti-diagonal, $(x, (x, x'))$.

Note that in each case A is a partially ordered set (there is at most one morphism between any two objects), and we illustrate the incidence relations between the objects in Fig. 1. The figure also illustrates how to construct functors $F : A \rightarrow \mathbf{Top}$, which are cofibrant with respect to the Reedy model structure on this functor category, where A itself is considered a directed Reedy category relative to the obvious dimension assignment (points of dimension zero, lines of dimension one, and the hemispheres of dimension two). We conclude that the homotopy colimit of F (in \mathbf{Top}) is weakly equivalent to the ordinary colimit of F , which is seen to be equivalent to $S^2 \vee S^1$, $S^2 \vee S^1 \vee S^1$, S^1 and $S^1 \vee S^1$, respectively for the cases (w, \cdot) , $(w, ')$, (we, \cdot) and $(we, ')$. Since F takes values in contractible spaces, this homotopy colimit represents in each case the homotopy type of the nerve of A . It follows that A is not aspheric.

It remains to see that the inclusion $A \hookrightarrow \mathbb{C}/\square^1 \times \square^1$ is in each case a weak equivalence. For this we use the dual of Quillen's Theorem A, and show that for each object $(x_1 \cdots x_n, (s, t))$ (written (s, t) for short) of $\mathbb{C}/\square^1 \times \square^1$, the coslice category $B_{(s, t)} = {}_{(s, t)}\backslash A$ has an initial object, and is hence aspheric.

As in Maltiniotis (*loc. cit.*), we make case distinctions on (s, t) . Note that for the languages under consideration, a term can have at most one free variable.

- $(s, t) = (x', y')$ for distinct variables x, y . The initial object is a hemisphere (the northern one in the presence of exchange, the southern one without exchange and in case x, y appear in reversed order in the variable context).
- $(s, t) = (x, x)$ or (x', x') . The initial object is then the diagonal in A .
- $(s, t) = (x, x')$ or (x', x) . The initial object is here the anti-diagonal.
- $(s, t) = (i, x')$ or (x', i) . The initial object is the corresponding side of the square.
- $(s, t) = (i, j)$. The initial object is then the corresponding vertex of the square. \square

We remark that this proof method fails in the presence of a connection, where terms can now refer to multiple variables. And in the presence of contraction, the diagonal (x, x) is incident on the northern hemisphere, which is then a maximal element in A , and hence A is contractible. These observations explain why it is only those four cases that can fail to give strict test categories.

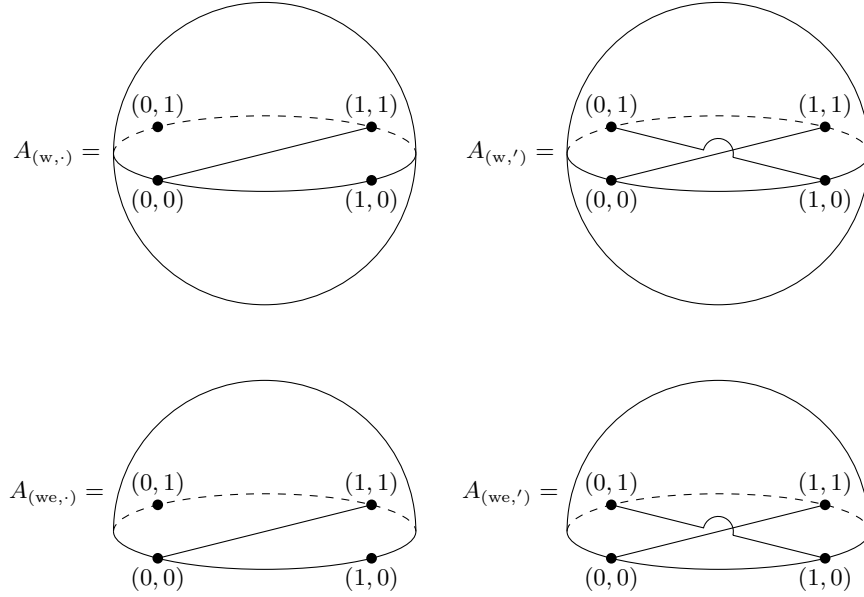


Fig. 1. The partially ordered sets $A = A_{(a,b)}$ in their topological realizations.

4 Conclusion

We have espoused a systematic algebraic framework for describing notions of cubical sets, and we have shown that all reasonable cubical sites are test categories. We have shown that a cubical site is a strict test category precisely when it is cartesian monoidal or includes one of the connections.

To improve our understanding of the homotopy theory of each cubical site, we need to investigate the induced Quillen equivalence between the corresponding category of cubical sets $\widehat{\mathbb{C}}$ with the Cisinski model structure and the category of simplicial sets $\widehat{\Delta}$ with the Kan model structure. For instance, we can ask whether the fibrations are those satisfying the cubical Kan filling conditions. Furthermore, to model type theory, one should probably require that the model structures lift to algebraic model structures. We leave all this for future work.

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